ADTs, Asymptotics II, BSTs

Exam-Level 06: October 7, 2024

1 Finish the Runtimes

Below we see the standard nested for loop, but with missing pieces!

```
for (int i = 1; i < _____; i = _____) {
    for (int j = 1; j < _____; j = _____) {
        System.out.println("Circle is the best TA");
}
</pre>
```

For each part below, **some** of the blanks will be filled in, and a desired runtime will be given. Fill in the remaining blanks to achieve the desired runtime! There may be more than one correct answer.

Hint: You may find Math.pow helpful.

```
(a) Desired runtime: \Theta(N^2)
    for (int i = 1; i < N; i = i + 1) {</pre>
        for (int j = 1; j < i; j = ____) {
            System.out.println("This is one is low key hard");
        }
   }
(b) Desired runtime: \Theta(\log(N))
    for (int i = 1; i < N; i = i * 2) {
        for (int j = 1; j < ____; j = j * 2) {
            System.out.println("This is one is mid key hard");
        }
   }
(c) Desired runtime: \Theta(2^N). \frac{2^N}{N} is a valid answer, could you think of another?
    for (int i = 1; i < N; i = i + 1) {
        for (int j = 1; j < ____; j = j + 1) {
            System.out.println("This is one is high key hard");
        }
   }
(d) Desired runtime: \Theta(N^3)
    for (int i = 1; i < ____; i = i * 2) {
        for (int j = 1; j < N * N; j = ____) {
            System.out.println("yikes");
        }
   }
```

2 Asymptotics is Fun!

(a) Using the function g defined below, what is the runtime of the following function calls? Write each answer in terms of N. Feel free to draw out the recursion tree if it helps.

```
void g(int N, int x) {
2
        if (N == 0) {
3
             return;
        }
4
        for (int i = 1; i <= x; i++) {
5
            g(N - 1, i);
6
7
        }
   }
8
   g(N, 1): \Theta(
   g(N, 2): \Theta(
```

(b) Suppose we change line 6 to g(N-1, x) and change the stopping condition in the for loop to $i \le f(x)$ where f returns a random number between 1 and x, inclusive. For the following function calls, find the tightest Ω and big O bounds. Feel free to draw out the recursion tree if it helps.

```
void g(int N, int x) {
if (N == 0) {
    return;
}

for (int i = 1; i <= f(x); i++) {
    g(N - 1, x);
}

g(N, 2): Ω( ), O( )</pre>
```

```
g(N, N): \Omega()
```

3 Is This a BST?

In this setup, assume a BST (Binary Search Tree) has a key (the value of the tree root represented as an int) and pointers to two other child BSTs, left and right.

(a) The following code should check if a given binary tree is a BST. However, for some trees, it returns the wrong answer. Give an example of a binary tree for which brokenIsBST fails.

```
public static boolean brokenIsBST(BST tree) {
    if (tree == null) {
        return true;
    } else if (tree.left != null && tree.left.key > tree.key) {
        return false;
    } else if (tree.right != null && tree.right.key < tree.key) {
        return false;
    } else {
        return brokenIsBST(tree.left) && brokenIsBST(tree.right);
    }
}</pre>
```

(b) Now, write isBST that fixes the error encountered in part (a).

Hint: You will find Integer.MIN_VALUE and Integer.MAX_VALUE helpful.

 $Hint\ 2$: You want to somehow store information about the keys from previous layers, not just the direct parent and children. How do you use the parameters given to do this?